

# The Stack and Procedures

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Computer Organization and Assembly Language

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## Overview

- The stack segment of a program is used for temporary storage of data and addresses.
- A *stack* is one-dimensional data structure.
- Items are added and removed from one end of the structure; that is, it is processed in a "last-in, first-out" manner.
- The most recent addition to the stack is called the top of the stack.

# Stack

- ▶ A program must set aside a block of memory to hold the stack.
- ▶ We declared a stack segment in the code as,
- ▶ `.STACK 100H`
- ▶ When the program is loaded in memory, SS will contain the segment number of the stack segment.
- ▶ For the preceding stack declaration, SP (stack pointer) is initialized to 100h, which represents empty stack.
- ▶ When the stack is not empty, SP contains the offset address of the top of the stack.

## *PUSH and PUSHF*

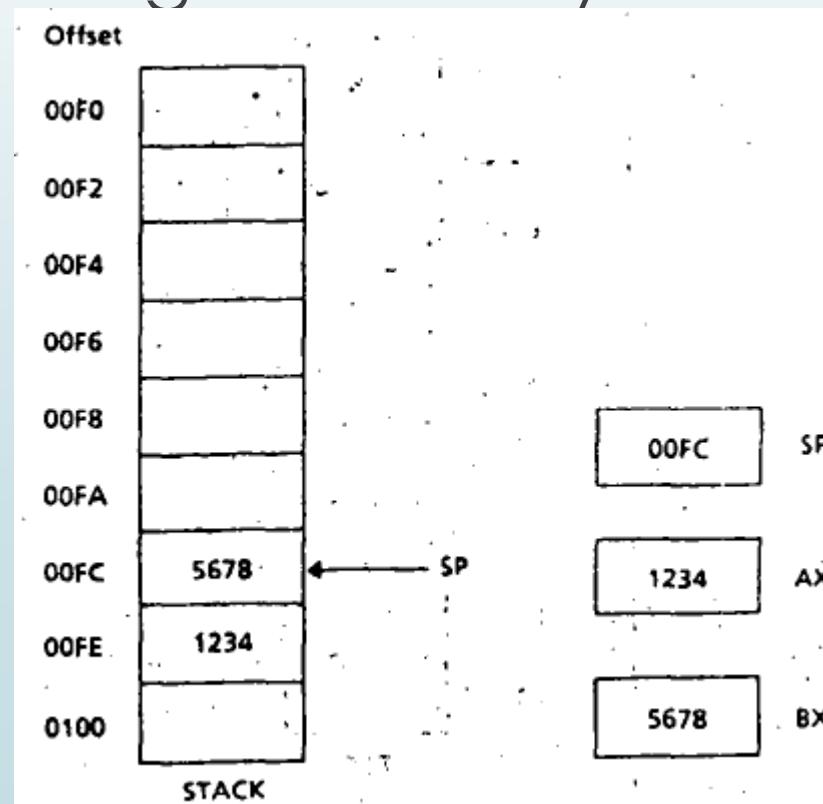
- ▶ To add a new data to the stack we PUSH it on stack as
- ▶ PUSH source
- ▶ where source is a 16-bit register or memory word. For example,
- ▶ PUSH AX
- ▶ PUSH execution does the following:
  - ▶ 1. SP is decreased by 2
  - ▶ 2. A copy of the source content is moved to the address specified by SS:SP.
- ▶ The instruction PUSHF, which has no operands, pushes the contents of the FLAGS register onto the stack.

## *SP Register*

- ▶ Initially, SP contains the offset address of the memory location immediately following the stack segment; the first PUSH decreases SP by 2, making it point to the last word in the stack segment.

## SP After PUSH Operations

- Since a PUSH decreases SP, the stack grows toward the beginning of memory.



## *POP and POPF*

- POP removes the top item from the stack as:
- POP destination
- where destination is a 16-bit register (except IP) or memory word. For example,
- POP BX
- Executing POP causes this to happen:
  - 1. The content of SS:SP (the top of the stack) is moved to the destination.
  - 2. SP is Increased by 2.

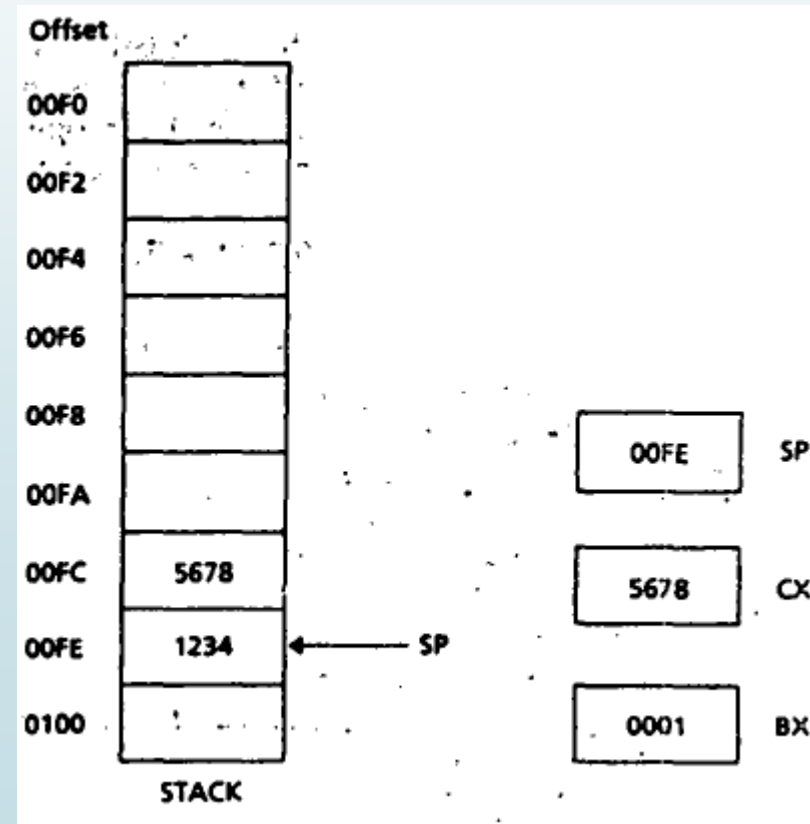


## POP and POPF

- ▶ The Instruction POPF pops the top of the stack into the FLAGS register.
- ▶ There is no effect of PUSH, PUSHF, POP, POPF on the flags.
- ▶ PUSH and POP are word operations, so a byte Instruction such as  
*PUSH DL*
- ▶ is illegal.
- ▶ So is a push of immediate data illegal in 8086, such as  
*PUSH 2*

## Stack After POP CX

- The figure below shows stack after *POP CX*



## OS and Stack

- The operating system also uses the stack.
- For example, to implement the INT 21h functions, DOS saves any registers it uses on the stack and restores them when the interrupt routine is completed.
- This does not cause a problem for the user since any values DOS pushes onto the stack are popped off by DOS before it returns control to the program.

# Procedures

- To solve a problem, it is decompose into subproblems that are easier to solve than the original problem.
- Procedures are used to solve these subproblems.
- One of the procedures is the main procedure, which is entry point to the program.
- The main procedure can call other procedures.
- The procedures can also call each other, or call itself (recursion).

## Procedures

- ▶ When one procedure calls another, control transfers to the called procedure and its instructions are executed;
- ▶ The called procedure returns control to the caller at the next instruction after the call statement.

## Procedure Syntax

- The syntax of procedure declaration is:

*name PROC type*

*;body of the procedure*

*RET*

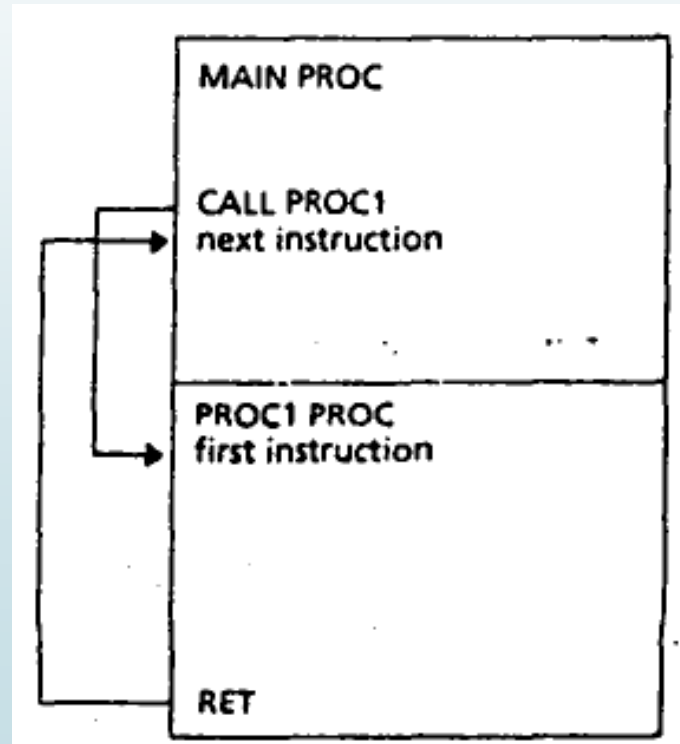
*name ENDP*

## Procedure Syntax

- Name is the user-defined name of the procedure.
- The optional **type** is NEAR or FAR (if omitted, NEAR is assumed).
- NEAR means that the statement that calls the procedure is in the same segment as the procedure itself.
- FAR means that the calling statement is in a different segment.

# Procedure Call

- The operating system also uses the stack.





## *RET Statement*

- ▶ The **RET** (return) instruction causes control to transfer back to the calling procedure.
- ▶ Every procedure (except the main procedure) should have a RET, usually in the end.

## *Passing Data Between Procedures*

- ▶ A procedure must have a way to receive values from the procedure that calls it, and a way to return results.
- ▶ Unlike high-level language procedures, assembly language procedures do not have parameter lists, so Registers and the Stack can be used for parameters and return values.

## CALL and RET

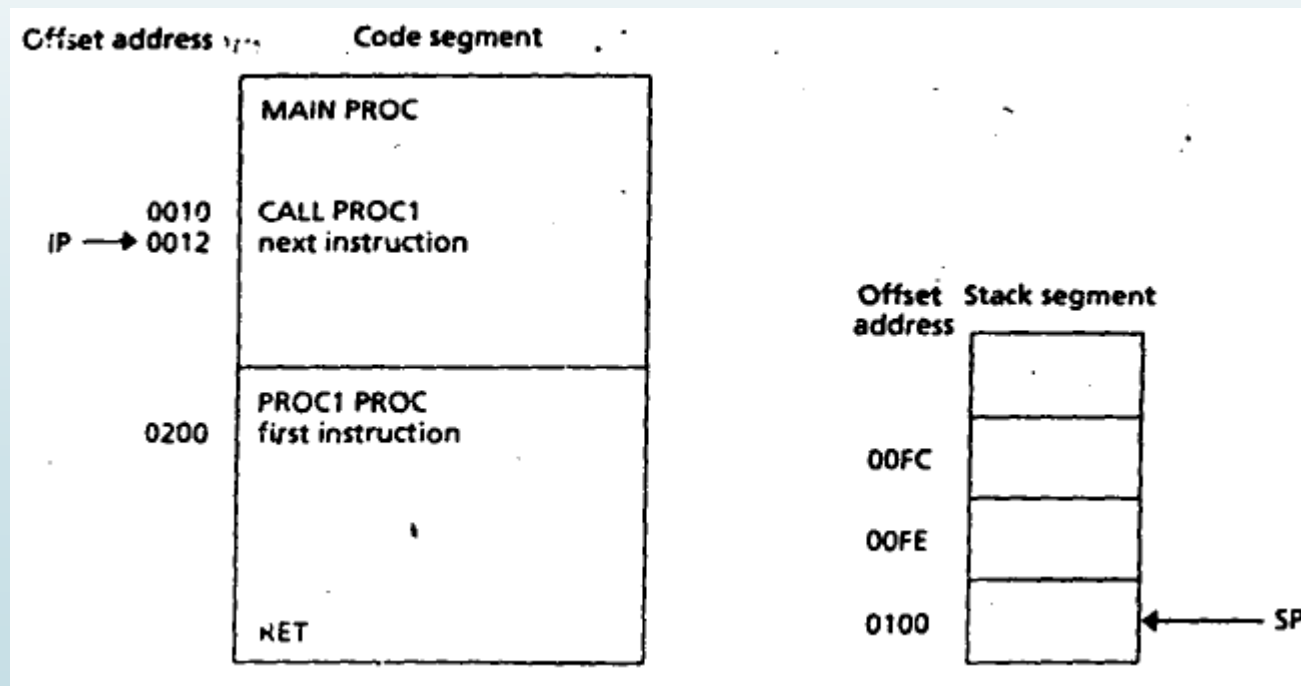
- To invoke a procedure, **CALL** instruction is used.
- There are two kinds of procedure calls, direct and indirect.
- The syntax of a direct procedure call is  
**CALL name**
- where name is the name of a procedure.
- The syntax of an indirect call is  
CALL address\_expression
- where address\_expression specifies a register or memory location containing address of a procedure.

## Executing a CALL instruction

- The return address to the calling program is saved on the stack.
- This is the offset of the next instruction after the CALL statement.
- The segment:offset of this instruction is in CS:IP at the time the call is executed.
- IP register gets the offset address of the first instruction of the procedure.
  - This transfers control to the procedure.

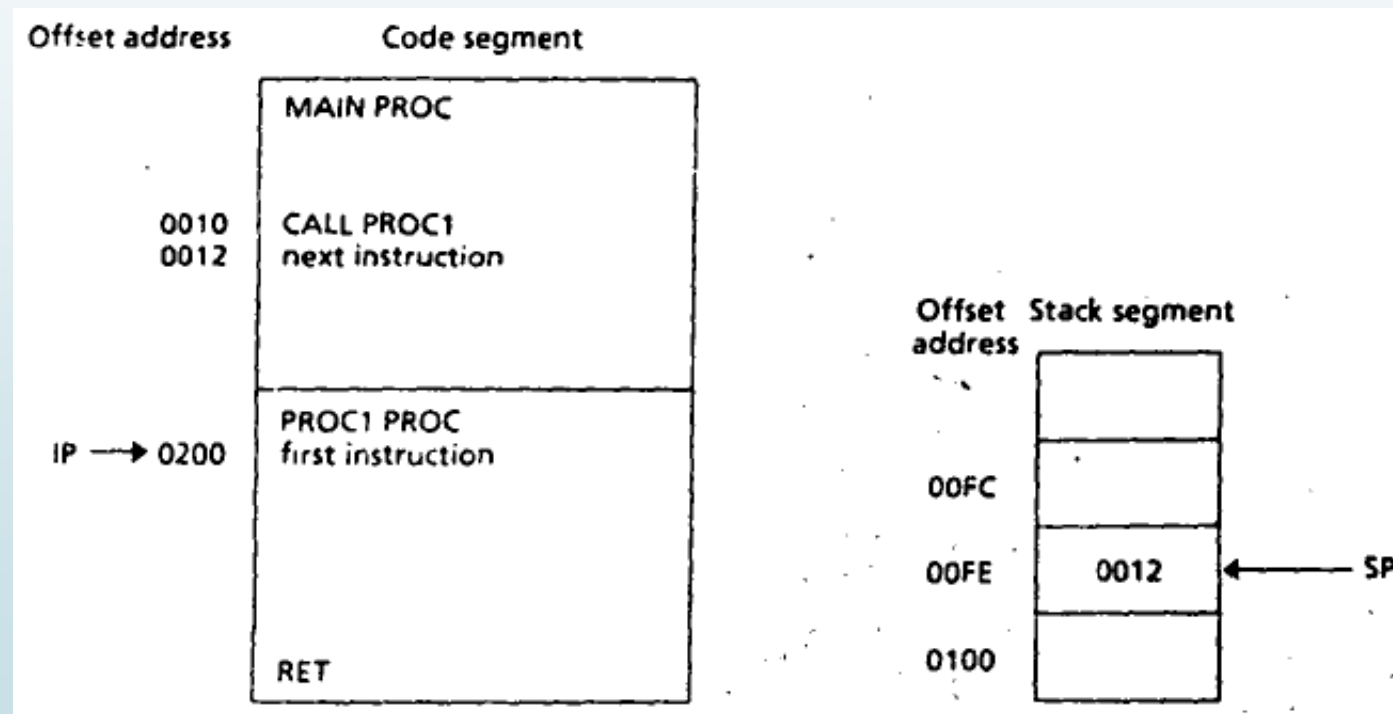
# Procedure Call and the Stack

- IP Register and the Stack before procedure call.



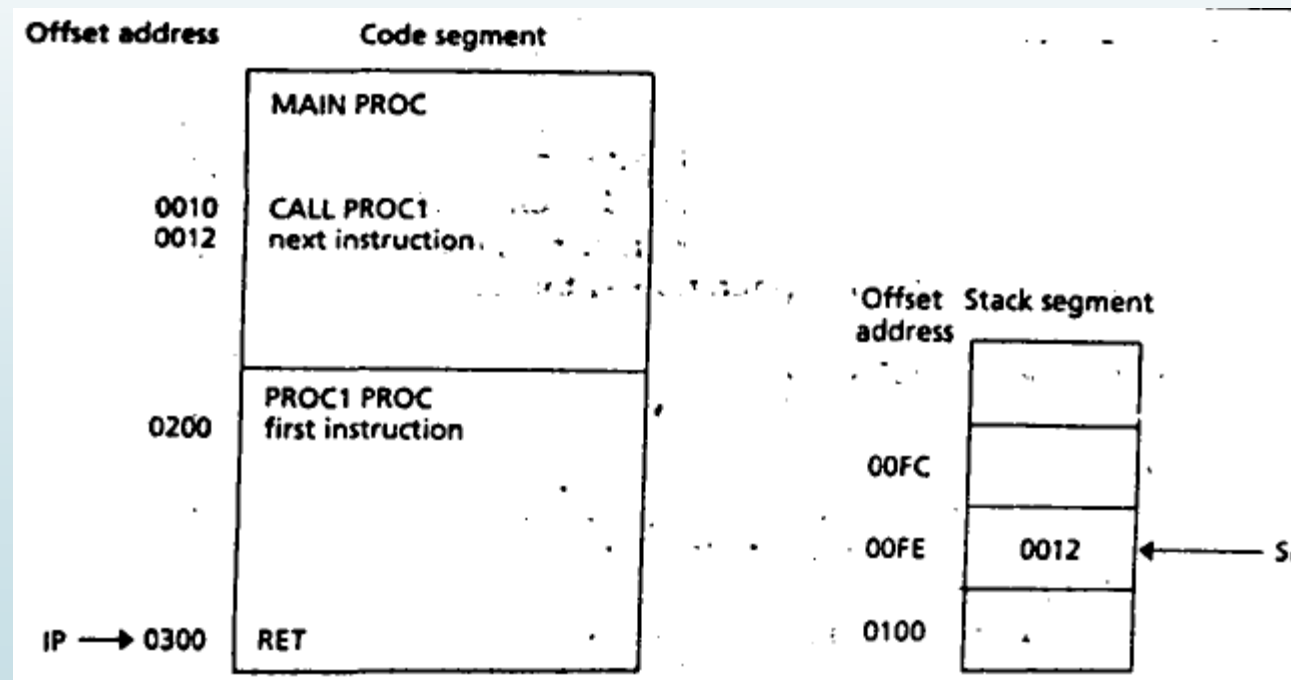
# Procedure Call and the Stack

- IP Register and the Stack after procedure call.



# Procedure Call and the Stack

- IP Register and the Stack before RET statement.



# Procedure Call and the Stack

- IP Register and the Stack after RET statement.

